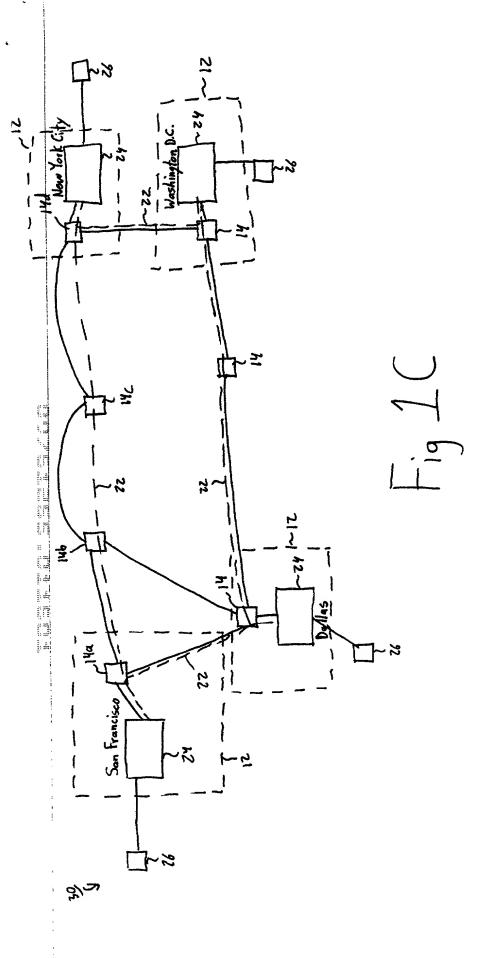


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Fig. 18

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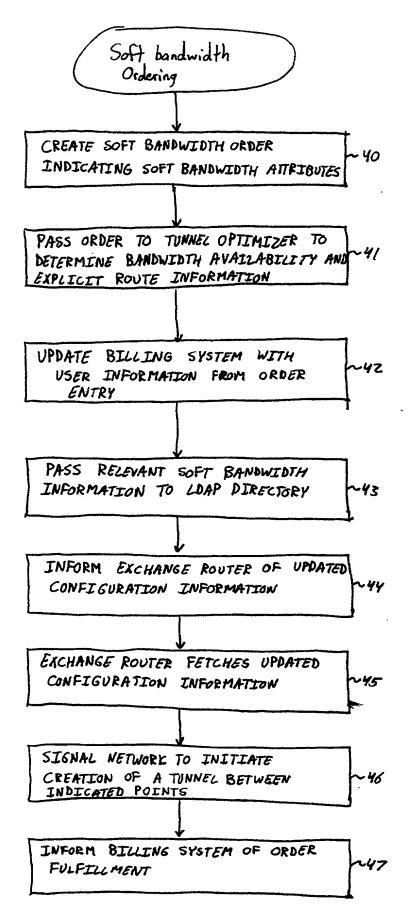
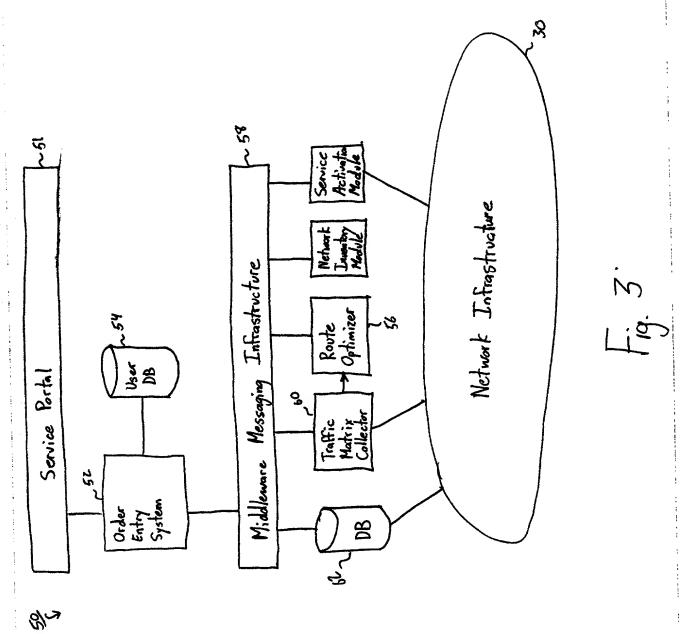


Fig. Z



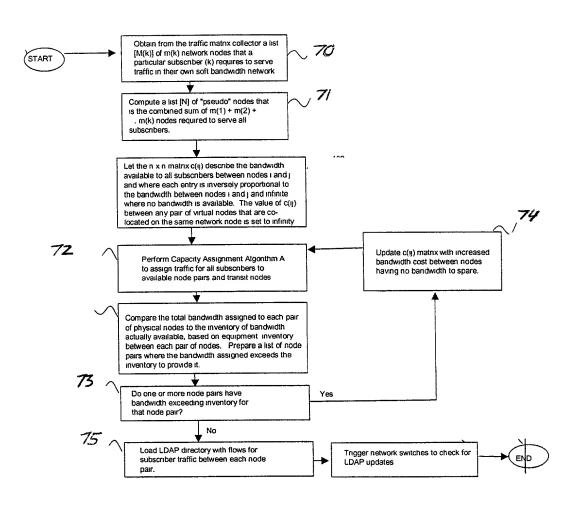


Fig 4

Input:

A list [N] of n network nodes having adequate equipment inventory to serve originating traffic and via traffic.

An $n \times n$ matrix [c(ij)] where each entry in the matrix is inversely proportional to the service bandwidth available between nodes i and j. Where there are no communications facilities directly connecting nodes i and j, the cost c(ij) is set to infinite.

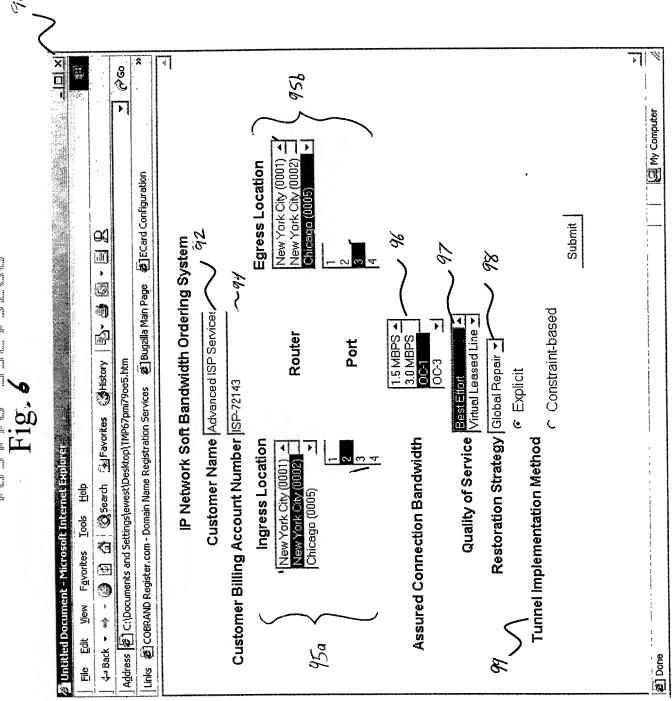
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Algorithm:
(Floyd-Warshall, 1962)
```

Output:

The route that a particular demand between any two points i and j may be found by looking up intermediate transit nodes found as values at the intersection of row i and column j in the $n \times n$ matrix e[ij].

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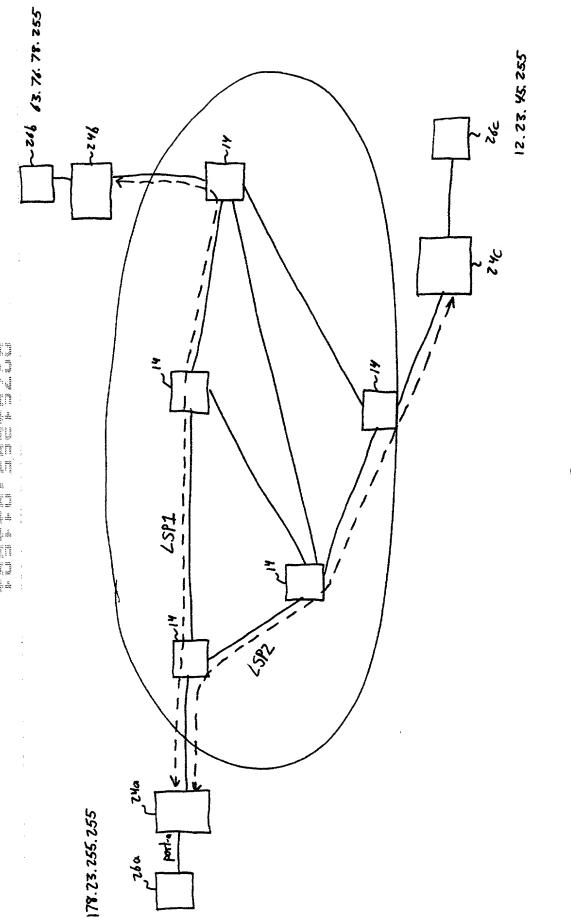
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